Trojan-tolerant Hardware + Supply Chain Security in Practice

Vasilios Mavroudis
Doctoral Researcher, UCL

Dan Cvrcek *CEO, Enigma Bridge*

Highlights

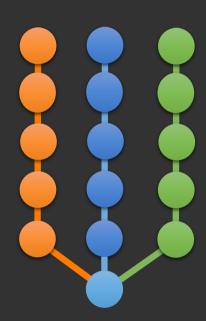
- The private life of keys
- Weak links of the supply chain
- Lessons learned from airplanes
- Demo of our crypto hardware
- Protocols, Maths & Magic
- Politics, Distrust & Hardware Security

The Private Life of Keys

- 1. Someone designs an integrated circuit (IC)
- 2. IC is fabricated
- з. IC is delivered to hardware vendor
- 4. Vendor loads firmware & assembles device
- 5. Device is sent to customer
- 6. Customer generates and stores key on the device

The Private Life of Keys

- 1. Someone designs the processor chip
- 2. Foundry fabricates the chip
- 3. Haulage transports chips
- 4. System vendor programs firmware
- 5. Distributors deliver a device to you
- 6. You create and use your key on the device



We can't protect all the steps

... but we can duplicate them

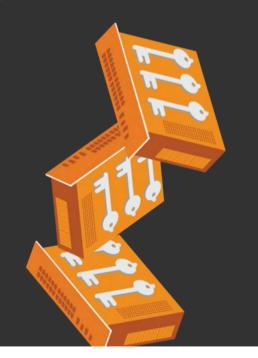
Hardware Security Modules

Physical computing device that safeguards and manages digital keys for strong authentication and provides cryptoprocessing.

Features:

- Cryptographic key generation, storage, management
- Tamper-evidence, Tamper-resistance, Tamper-response
- Security Validation & Certification

Crypto Operations are carried out in the device No need to output the private keys!



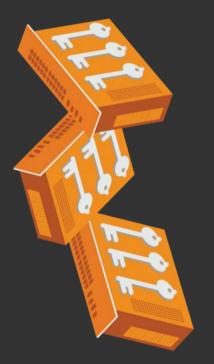
Hardware Security Modules

Common Applications

- Public Key Infrastructures
- Payment ProcessingSystems
- SSL Connections
- DNSSEC
- Transparent Data Encryption

Cost

- Hardware (>\$10k)
- Integration Cost
- Operational/Support



HSM Guarantees

- 1. Someone designs an integrated circuit (IC)
- 2. IC is fabricated
- 3. IC is delivered to hardware vendor
- 4. Vendor loads firmware & assembles device
- 5. Device is sent to customer
- 6. Customer generates and stores key on the device

What could go wrong?

Bugs

CVE-2015-5464

The HSM allows remote authenticated users to bypass intended key-export restrictions ...

Backdoors/HT?

THIS 'DEMONICALLY CLEVER' BACKDOOR HIDES IN A TINY SLICE OF A COMPUTER CHIP

NSA's Own Hardware Backdoors May Still Be a "Problem from Hell" Expert Says NSA Have Backdoors Built Into Intel And AMD Processors

Snowden: The NSA planted backdoors in Cisco products



Proposed Solutions

- Trusted Foundries
 - Very expensive
 - Prone to errors/bugs
- Split-Manufacturing
 - Still Expensive
 - □ Prone to errors/bugs

- Post-fabrication Inspection
 - Expensive (+ re-tooling)
 - A huge pain, doesn't scale

Arms Race

- Adversaries always one step forward
- Can never be 100% certain

It's safe to assume we will never win

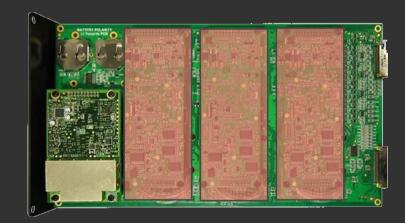


Solution from the sky and space

Lockstep systems are fault-tolerant computer systems that run the same set of operations at the same time in parallel.

- Dual redundancy allows error detection and error correction
- Triple redundancy automatic error correction, via majority vote
 → Triple Redundant 777 Primary Flight Computer





Safety != Security

Fault-tolerant systems are built for safety and the computations are simply replicated

Not enough for security!

Keys would have to be copied across all processors

Security of our keys would depend on the weakest link

Our Solution

- 1. Someone designs an integrated circuit (IC)
- 2. IC is fabricated
- 3. IC is delivered to hardware vendor
- 4. Vendor loads firmware & assembles device
- 5. Device is sent to customer
- 6. Customer generates and stores key on the device

Who we are

Vasilios
Mavroudis
Doctoral Researcher,

UCL

Dan Cvrcek
CEO, Enigma Bridge

George Danezis

Professor, UCL

Petr Svenda

Assistant Professor, MUni CTO, Enigma Bridge

Ingredients of the Solution

- 1. Hardware Components (IC) 2. Cryptographic Protocols
 - Independent Fabrication
 - Non-overlapping Supply Chains
 - Programmable
 - □ Affordable
 - Bonus if COTS

- - No single trusted party
 - □ Full Distribution of Secrets
 - Distributed Processing
 - Provably Secure (i.e., Math)

Smart Cards

Many Independent Manufacturers

- Private Fabrication Facilities
- Disjoint Supply Chains (location, factories, design)

Programmable Secure Execution Environment

- □ NIST FIPS140-2 standard, Level 4
- □ Common Criteria EAL4+/5+

Off-the-shelf Cost \$5-\$40

Multiparty Computation Protocols

Distributed Operations

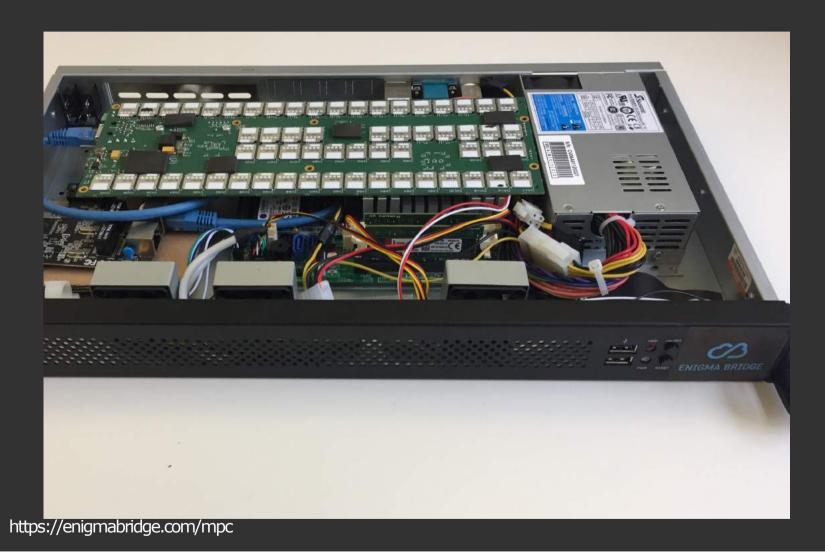
- □ Random number Generation
- □ Key Pair Generation
- □ Decryption
- □ Signing

Provably Protect against

- □ *N-1* Malicious & Colluding parties
- N Malicious & non-colluding parties



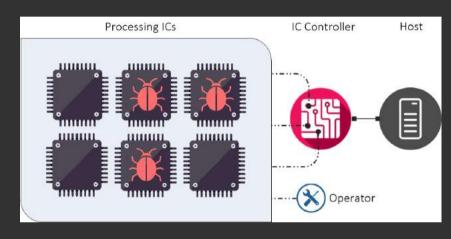
THE prototype

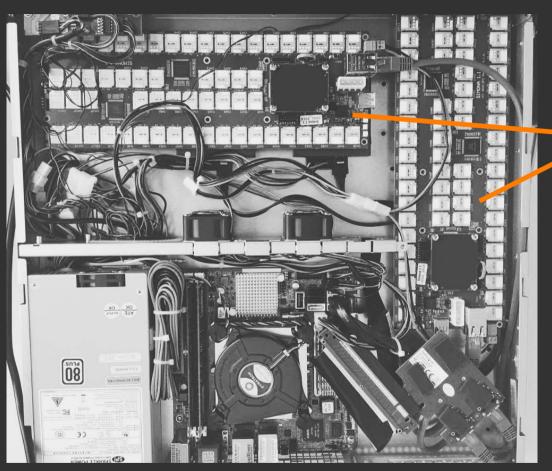


Many Smart Cards

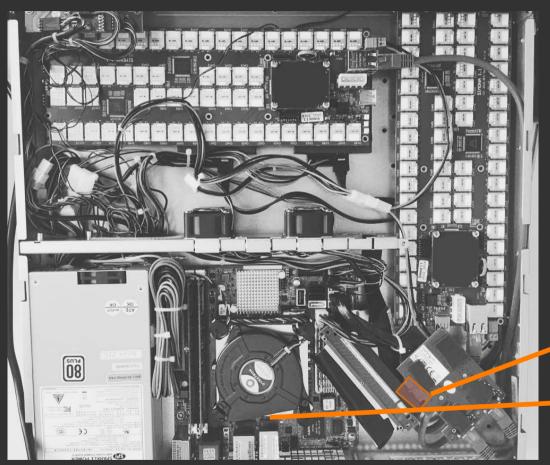
Components

- Multiples of 120 smartcards
- TCP/UDP access to smartcards
- FPGA manages the communication bus
 - 1Gbit/s bandwidth for requests

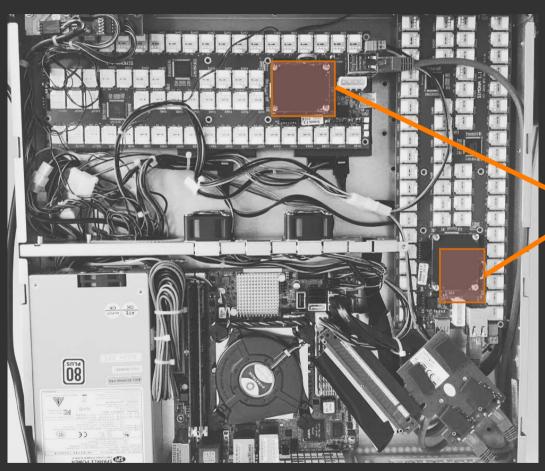




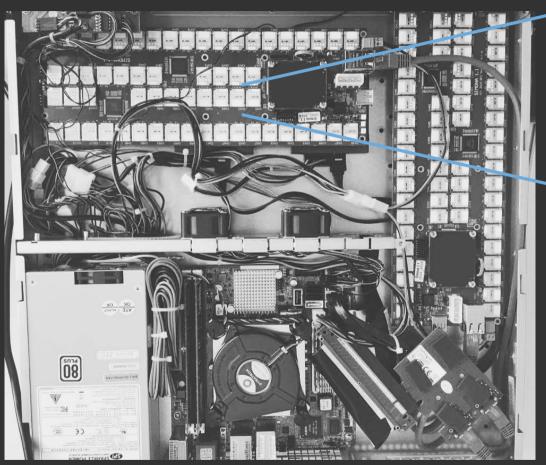
SIMONA boards with 120 JCs

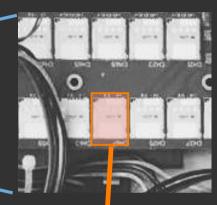


Gigabit link to untrusted Linux server



FPGA - JC /ISO7816->TCP

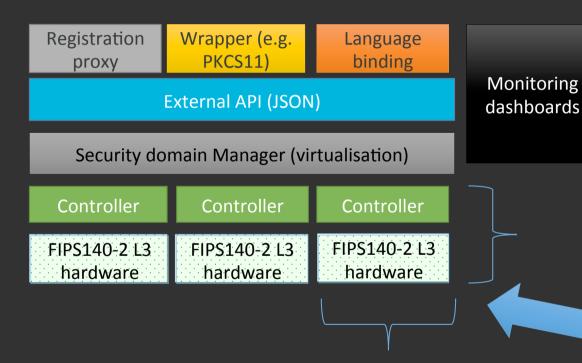


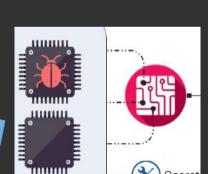


JavaCards

- FIPS140-2 Level 3
- CC EAL5+

Plugging it into a cloud service





Giving smart-cards an infrastructure

mpc.enigmabridge.com

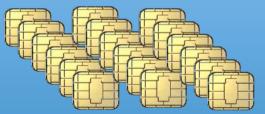
MacBook-2



MPC RESTful server

SIMONA board (192.168.42.10)

LAS VEGAS



nttps://enigmapnage.com/mpc

SIMONA board (84.92.209.143)

CAMBRIDGE UK



Demo 1

Giving smart-cards an infrastructure

mpc.enigmabridge.com

MacBook-2



MPC RESTful server

SIMONA board (192.168.42.10) **LAS VEGAS**

https://enigmabridge.com/mpc

SIMONA board (84.92.209.143)

CAMBRIDGE UK



- ARM

jCardSim

- Intel

- SPARC

Giving smart-cards an infrastructure

MPC RESTful server

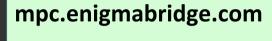
SIMONA board (192.168.42.10)
LAS VEGAS

MacBook-2



Demo 2

Visualizing Cryptography

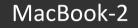


Node-red

- RESTful requests (switch evil)
- MPC key generation
- web-socket servers

MPC RESTful server





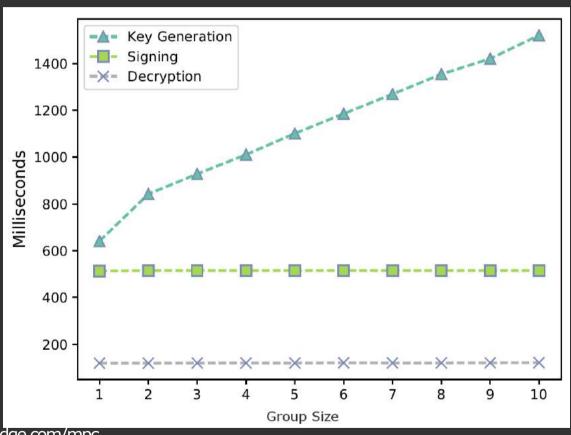


ICs with Hardware Trojans

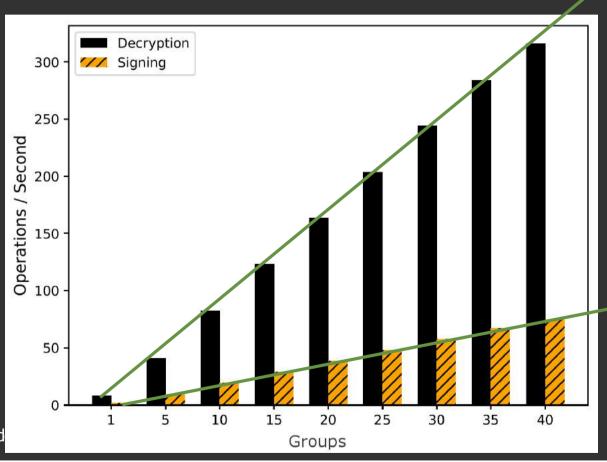
Demo 3

PERFORMANCE

Tolerance vs Runtime







Linear, no bottleneck

https://enigmabrid

Protocols

Key Points

- No single IC is trusted with a secret (e.g., private key)
- Misbehaving ICs can be detected by honest ones
- If one IC is excluded from any protocol, user can tell

Bonus: Minimize interaction between ICs for performance

Sharing a Secret

- Split a secret in shares
- The secret can be reconstructed later
- Without *sufficient* shares not a single bit is leaked
- Splitting Parameters:
 - How many shares the secret is split into (n)
 - How many shares you need to reconstruct the secret (t)

In our case: Each 3 ICs hold shares for a secret



Classic Key Generation

Single IC System

- 1. Bob asks for new key pair
- Backdoored IC generates compromised key
- 3. Private Key is "securely" stored
- 4. Weak Public key is returned

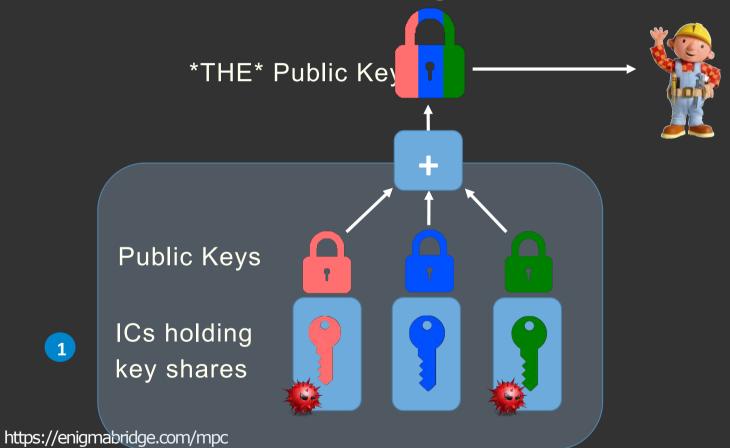
Problems

- Malicious IC has full access to the private key
- Bob can't tell if he got a "bad" key

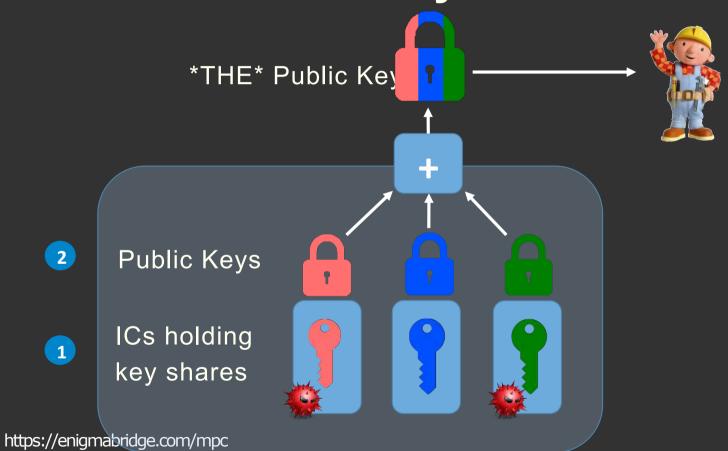


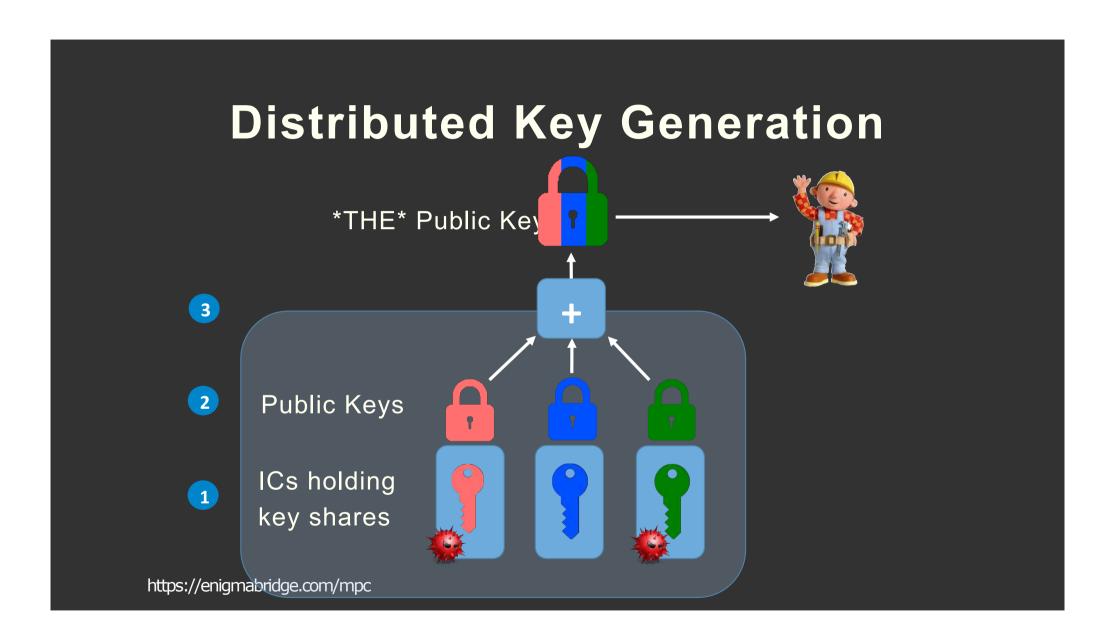


Distributed Key Generation

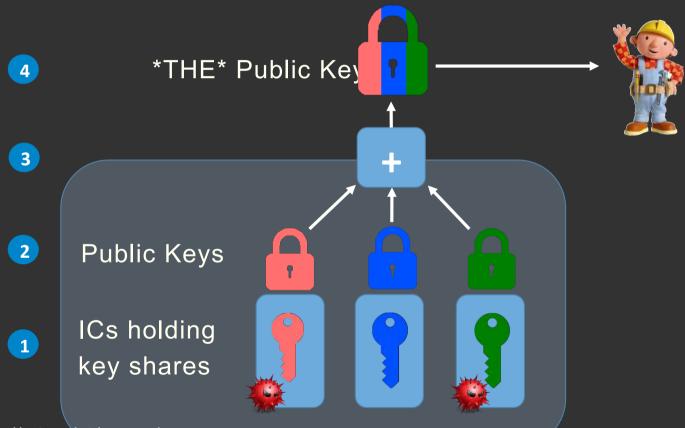








Distributed Key Generation



Distributed Key Generation

Key Points

- No single IC is trusted with a secret (e.g., private key) ✓
- Misbehaving ICs can be detected by honest ones 🗸
- If one IC is excluded from any protocol, user can tell 🗸

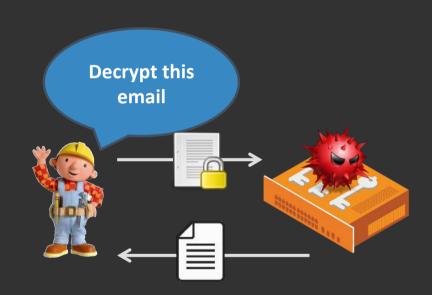
Bonus: Minimize interaction between ICs for performance 🗶

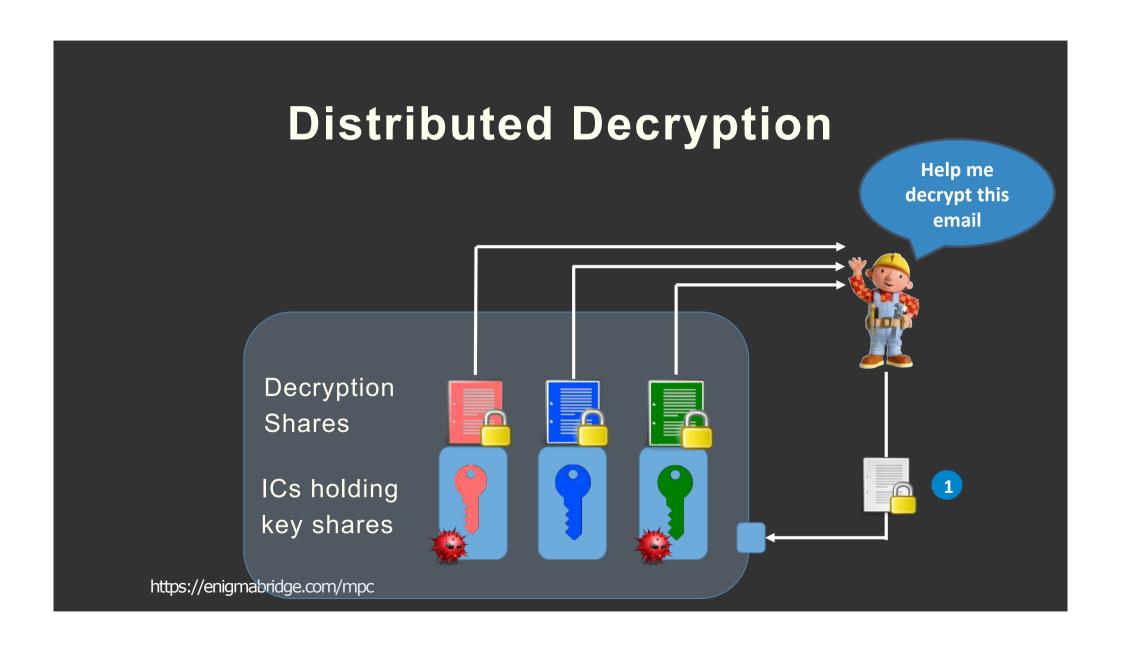
Classic Decryption

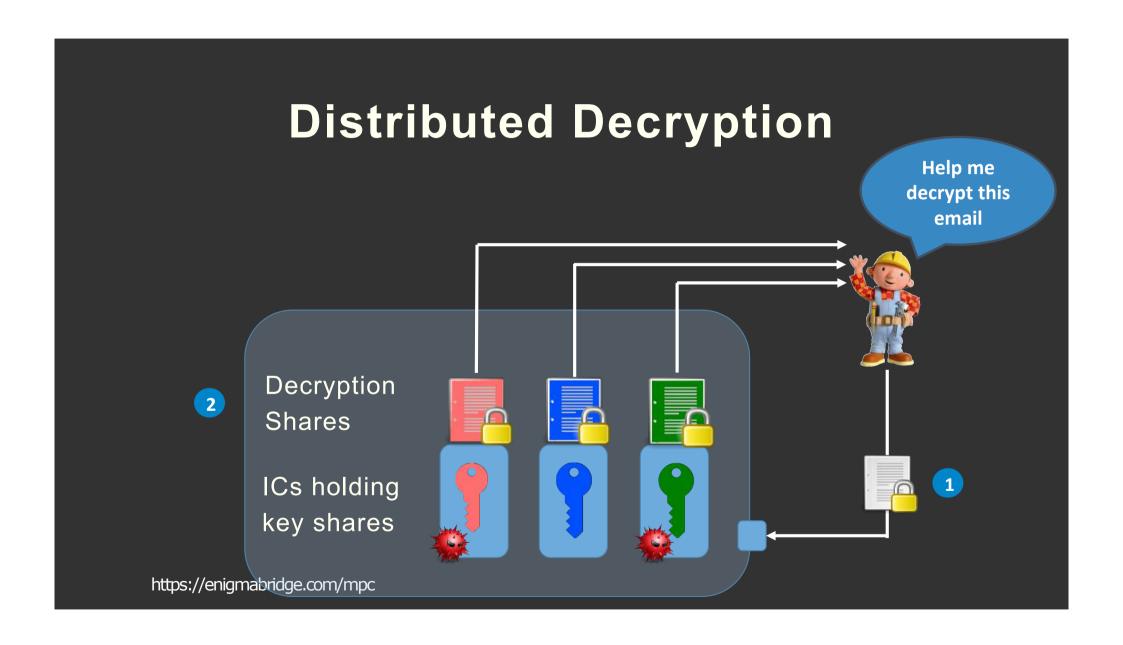
Single IC System

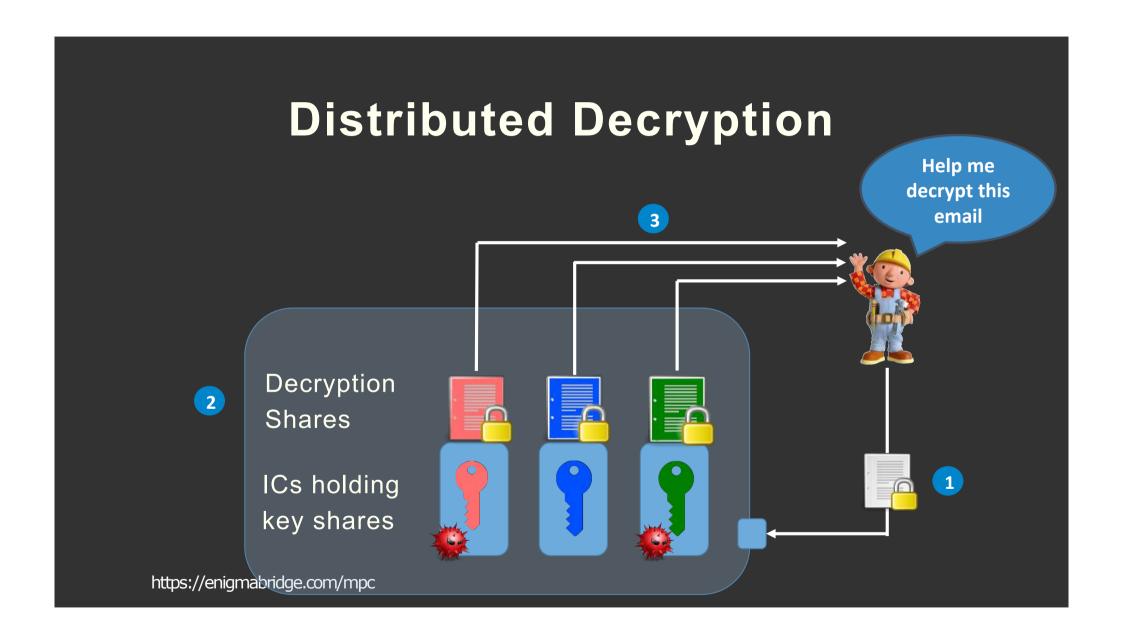
- 1. Bob asks for ciphertext decryption
- 2. Backdoored IC decrypts ciphertext
- 3. Bob retrieves plaintext

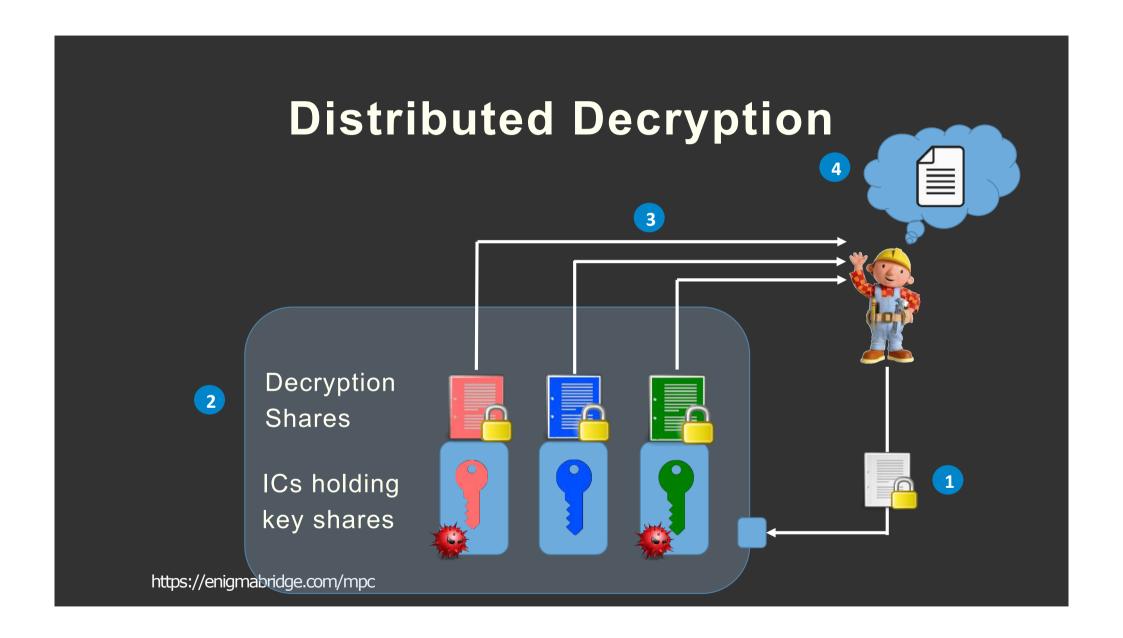
The IC needs full access to the private key to be able to decrypt ciphertexts.











Distributed Decryption

Key Points

- No single IC is trusted with a secret (e.g., private key) ✓
- Misbehaving ICs can be detected by honest ones -
- If one IC is excluded from any protocol, user can tell 🗸

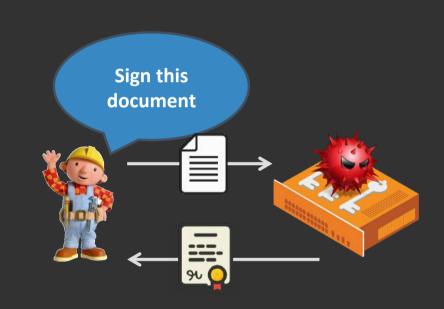
Bonus: Minimize interaction between ICs for performance 🗸

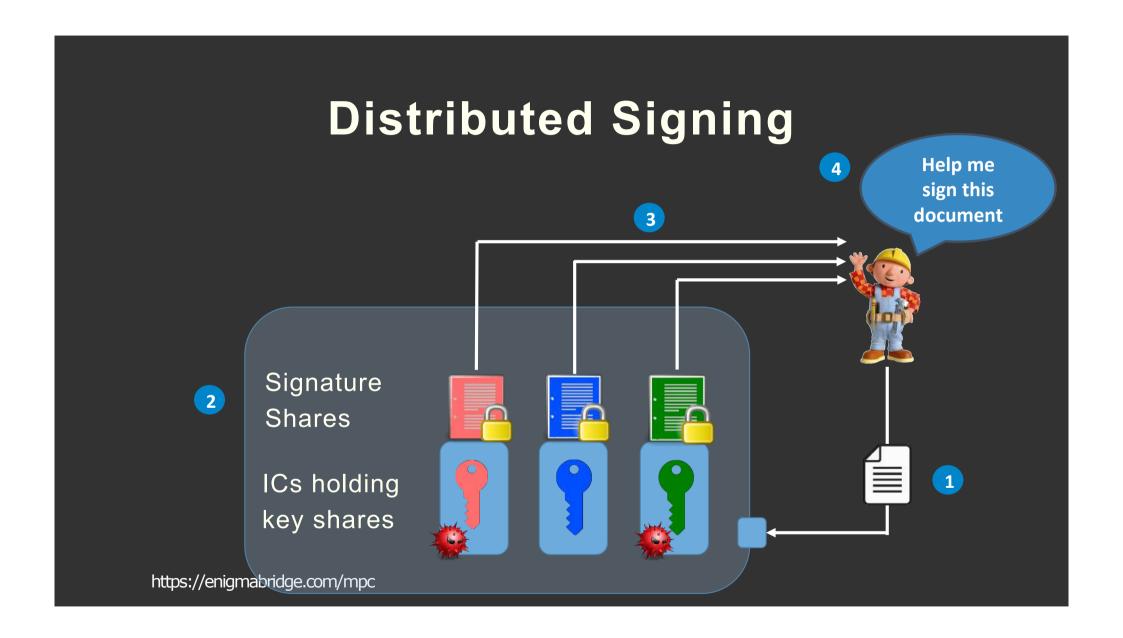
Classic Signing

Single IC System

- 1. Bob asks for document signing
- 2. Backdoored IC signs the plaintext
- 3. Bob retrieves signature

The IC needs full access to the private key to be able to sign plaintexts.





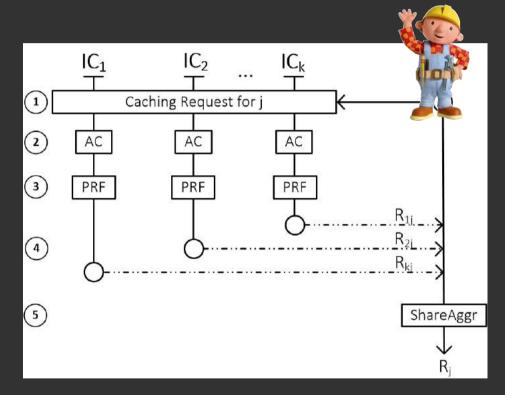
Distributed Signing I

Caching

- 1. Bob sends a caching request
- 2. The ICs verify Bob's authorization
- 3. Generate a random group element based on j
- 4. Bob sums the random elements

Properties

- Caching for thousands of rounds (j)
- Bob stores R https://enigmabridge.com/mpc



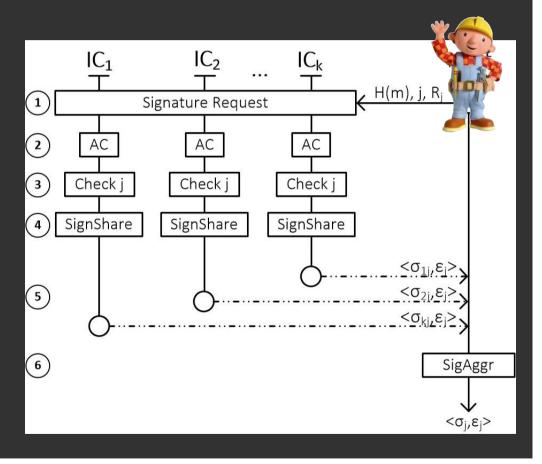
Distributed Signing II

Signing

- Bob asks for document signing & sends R_{i.} j, and the hash of m
- 2. ICs verify his authorization
- 3. ICs check if j has been used again
- 4. ICs compute their signature share
- 5. Bob sums all signature shares

Properties

- All ICs must participate
- Signation Signation Signation Signation Signation Signation Signature Sign



Distributed Signing

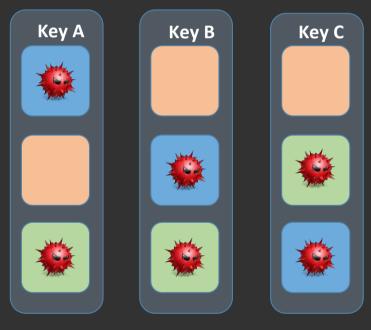
Key Points

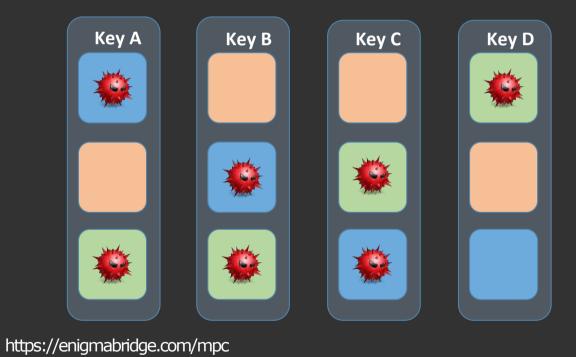
- No single IC is trusted with a secret (e.g., private key) ✓
- Misbehaving ICs can be detected by honest ones 🗸
- If one IC is excluded from any protocol, user can tell 🗸

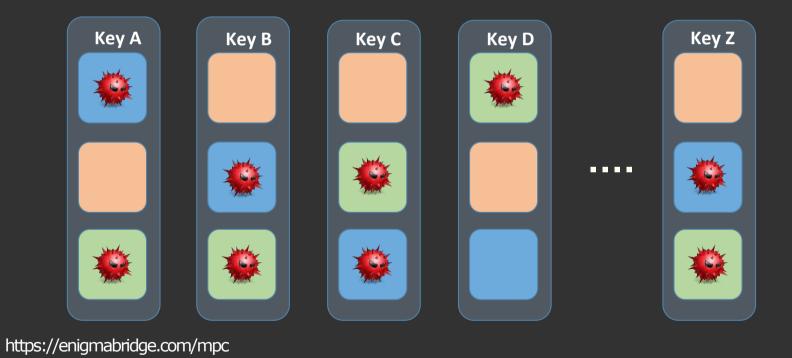
Bonus: Minimize interaction between ICs for performance 🗸



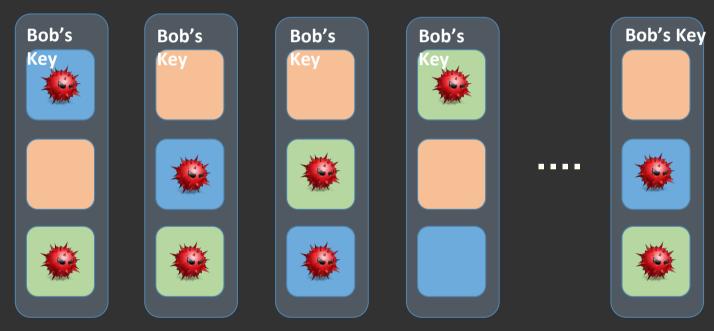




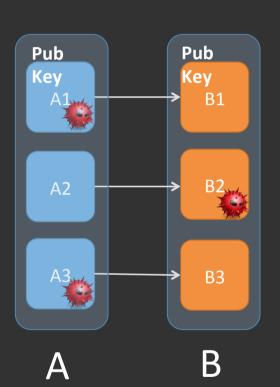




But how can all these groups have shares for the same key?

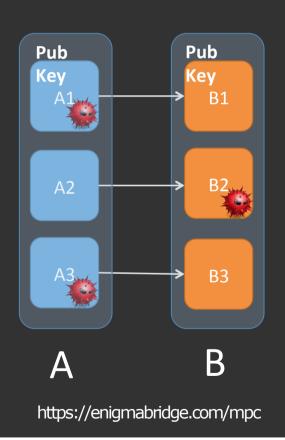


Key Replication



- 1. Group A generates a public key
- 2. A1, A2, A3 send their shares to B1, B2, B3
- 3. Each IC in B receives shares from A1, A2, A3
- 4. Each IC in B combines the 3 shares and retrieves its private key

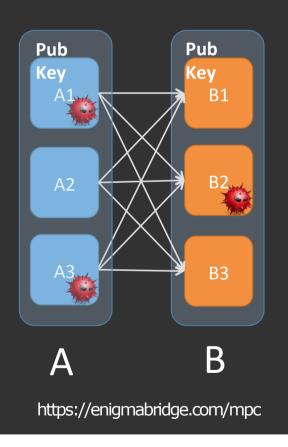
Key Replication



- 1. Group A generates a public key
- 2. A1, A2, A3 send their shares to B1, B2, B3
- 3. Each IC in B receives shares from A1, A2, A3
- 4. Each IC in B combines the 3 shares and retrieves its private key
- 5. A1, A3 and B2 collude

The adversary retrieves the secret!

Key Replication



- 1. Group A generates a public key
- 2. Then each IC in A splits its private key in three shares and sends them to B1, B2, B3
- 3. Each IC in B receives shares from A1, A2, A3
- 4. Each IC in B combines the 3 shares and retrieves its private key share

The full public keys of A and B are the same!

geopolitics

"We can guarantee security if there is at least one honest IC that is not backdoored or faulty."

"We can guarantee security if there is at least one honest IC that is not backdoored or faulty."

What if all ICs are malicious?

Government-level adversaries

- Deep access to fabrication facilities
- Very sophisticated techniques
- Very hard to detect their Backdoors/ Trojans
- Very secretive; highly classified
- Won't share their backdoor details

Government-level adversaries

- Deep access to fabrication facilities
- Very sophisticated techniques
- Very hard to detect their Backdoors/Trojans
- Very secretive; highly classified
- Won't share their backdoor details
- Unlikely to collude with anyone

"We can guarantee security even when all ICs are malicious, if at least one does not collude."



Conclusions & Future

New crypto hardware architecture

- For the first time, tolerates faulty & malicious hw
- Decent Performance
- Scales nicely; just keep adding ICs
- Suitable for commercial-off-the-shelf components
- Existing malicious insertion countermeasures are very welcome!

Trojan-tolerant Hardware + Supply Chain Security in Practice

find more at https://enigmabridge.com/mpc

Vasilios Mavroudis
Doctoral Researcher, UCL

Dan Cvrcek *CEO, Enigma Bridge*